

## **Lab #2**

# **Static Random-Access Memory**

## **(SRAM)**

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**Objective:**

The goal of this lab is to create a static random-access memory chip using D-Latches, also known as SRAM. An SRAM is used to store multiple bits at various address locations. The SRAM that we will be designing will hold 16-bit wide values and could display in the 7-segment HEX displays and decimals. We will be creating 2 forms of SRAM, one 16 x 4 SRAM and one 16 x 32 SRAM. The 16 x 4 SRAM will contain 16 locations, 4 bits each and the 16 x 32 SRAM will contain 16 locations, 32 bits each. With the 16 x 32 SRAM, we will also learn how to design a way to input 32 bits to memory word using 4 times 8 switches and keys.

## Functionality and Specification:

### SR-Latch:

A SR-Latch is a circuit that compares two inputs and lets out a result of two outputs. It acts as two functions, set and reset, which are controlled by the inputs, “S” and “R”, respectively. When “R” is 0 and “S” is 0, then the output remains the same as the previous state. When “R” is 1 and “S” is 0, then the output resets the device, making “Q” become 0. When “R” is 0 and “S” is 1, then the output sets the device making “Q” become 1. When “R” is 1 and “S” is 1, then the output is undefined because you can't have it both set and reset at the same time, so it will be impossible to predict the output.

The truth table, and circuit are shown below:

S	R	Q	Q'	State
0	0	-	-	No Change
0	1	0	1	Reset
1	0	1	0	Set
1	1	?	?	Undefined

Table 1: Truth Table for SR-Latch

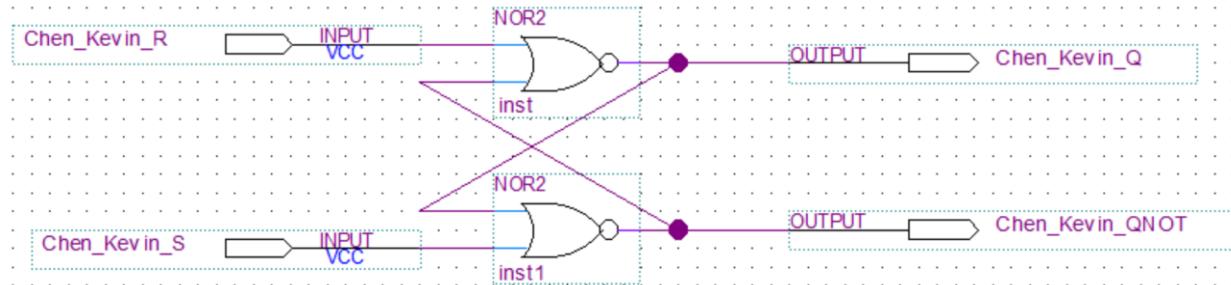


Figure 1: Block Diagram/Schematic File of SR-Latch (Chen\_Kevin\_SR-Latch.BDF)

### Control SR-Latch:

A Control SR-Latch is a circuit that compares three inputs and lets out a result of two outputs. It acts the same as a SR-Latch except with an additional input, Control (C). The control input basically determines if the latch should change its state and take in the current inputs of “S” and “R”. When “R” is 0 and “S” is 0, then the output remains the same as the previous state. When “C” is 0, then it doesn’t matter what “R” and “S” is because the output will always remain the same as the previous state. When “C” is 1, then the inputs and output will act the same way as the SR-Latch.

The truth table, and circuit are shown below:

C	S	R	Q <sub>n</sub>	Q <sub>n'</sub>	State
0	X	X	Q	Q'	No Change
1	0	0	Q	Q'	No Change
1	0	1	0	1	Reset
1	1	0	1	0	Set
1	1	1	?	?	Undefined

Table 2: Truth Table for Control SR-Latch

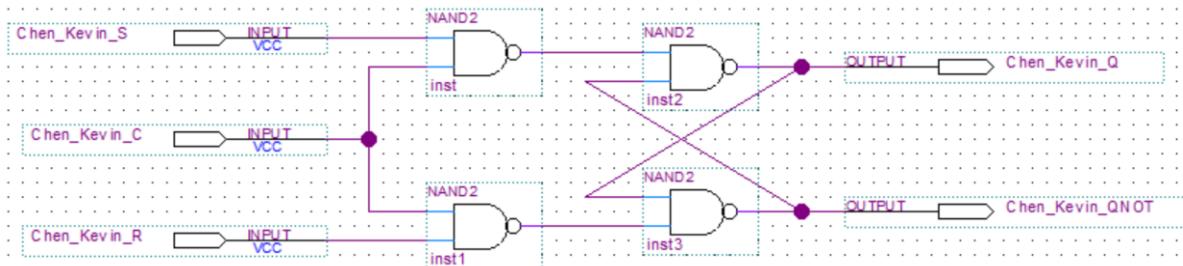


Figure 2: Block Diagram/Schematic File of Control SR-Latch

(Chen\_Kevin\_Control\_SR-Latch.BDF)

### D-Latch:

A D-Latch is a circuit that compares two inputs and lets out a result of two outputs. It acts the same as a Control SR-Latch except with a combined input, “D”. The “D” input basically combines the “S” and “R” input into one input. When “C” is 0, then it doesn’t matter what “D” is because the output will always remain the same as the previous state. When “C” is 1, then the output would depend on what “D” is. If “D” is 0, then the output resets the device, making “Q” become 0. If “D” is 1, then the output sets the device, making “Q” become 1.

The truth table, and circuit are shown below:

C	D	$Q_n$	$Q'_n$	State
0	X	Q	$Q'$	No Change
1	0	0	1	Reset
1	1	1	0	Set

Table 3: Truth Table for D-Latch

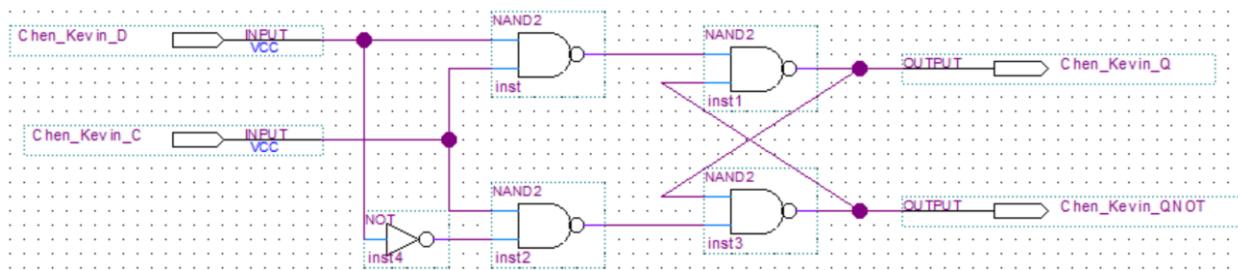


Figure 3: Block Diagram/Schematic File of D-Latch (Chen\_Kevin\_D-Latch.BDF)

SRAM:

A SRAM is a circuit that compares three inputs and lets out a result of one output. It's made from two D-Latches, which could also be described as a Master-Slave D Flip-Flop. A Master-Slave D Flip-Flop isn't the same a regular D-Latch because its state change does not depend on the state of the clock signal, but the transition of the clock signal from zero to one. A SRAM is also made of a Tristate Logic Buffer, which compares two inputs and lets out a result of one output. The two inputs are "A" and "B" and it only outputs "A" when "B" is 1, otherwise the output is cut off completely and no signal is sent. The three inputs of a SRAM are "IN", "Select\_Chip", and "Write\_Enable". "IN" is the latch data input. "Select\_Chip" is the input that activates the cell when it is high; when it's low, then the cell will not store data, and there will not be an output. "Write\_Enable" is the input that allows the latch to store the data from the "IN" data input.

The truth table, and circuit are shown below:

Clock	D <sub>M</sub>	C <sub>M</sub>	D <sub>S</sub>	C <sub>S</sub>	Q <sub>out</sub>
0	d1	1	q1	0	Q0
1	d2	0	q1	1	Q1
0	d3	1	q3	0	Q1
0	d4	1	q4	0	Q1
1	d5	0	q4	1	Q4
1	d6	0	q4	1	Q4

Table 4: Truth Table for Master-Slave D Flip-Flop

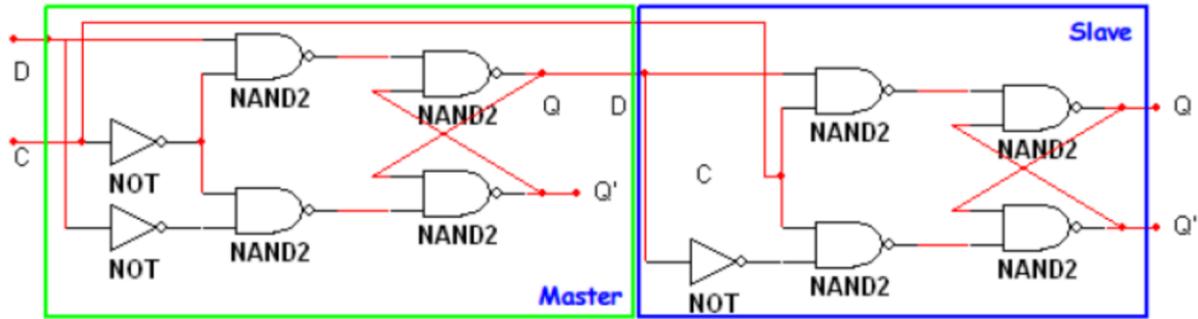


Figure 4: Master-Slave D Flip-Flop Design

A	B	C
0	1	0
1	1	1
X	0	Z

Table 5: Truth Table for Tristate Logic Buffer

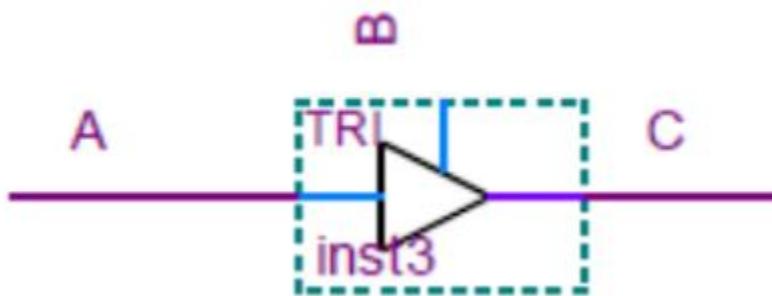


Figure 5: Tristate Logic Buffer Design

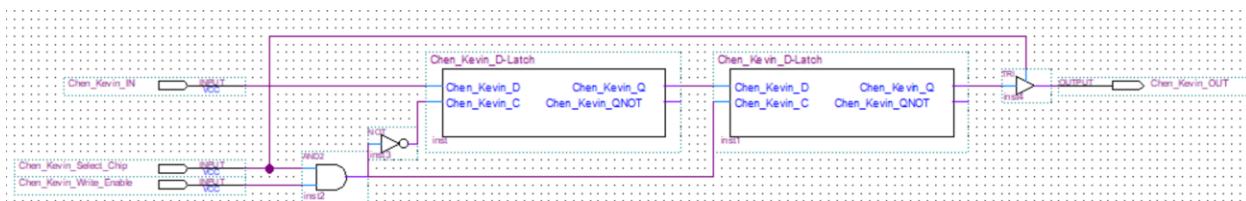


Figure 6: Block Diagram/Schematic File of SRAM (Chen\_Kevin\_SRAM.BDF)

### 16 x 1 SRAM:

A 16 x 1 SRAM is a circuit that's made of 16 SRAM that's all stacked into one file. It takes in the same input as a SRAM, except now, the data values are in a form of vectors. Dec[15..0] acts as "Select\_Chip" from SRAM, except being a vector from 0 to 15. "IN" acts as "IN" from SRAM, and "WR" acts as "Write\_Enable" from SRAM. The output of the 16 x 1 SRAM, "out", is the same as "OUT" from SRAM.

The circuit is shown below:

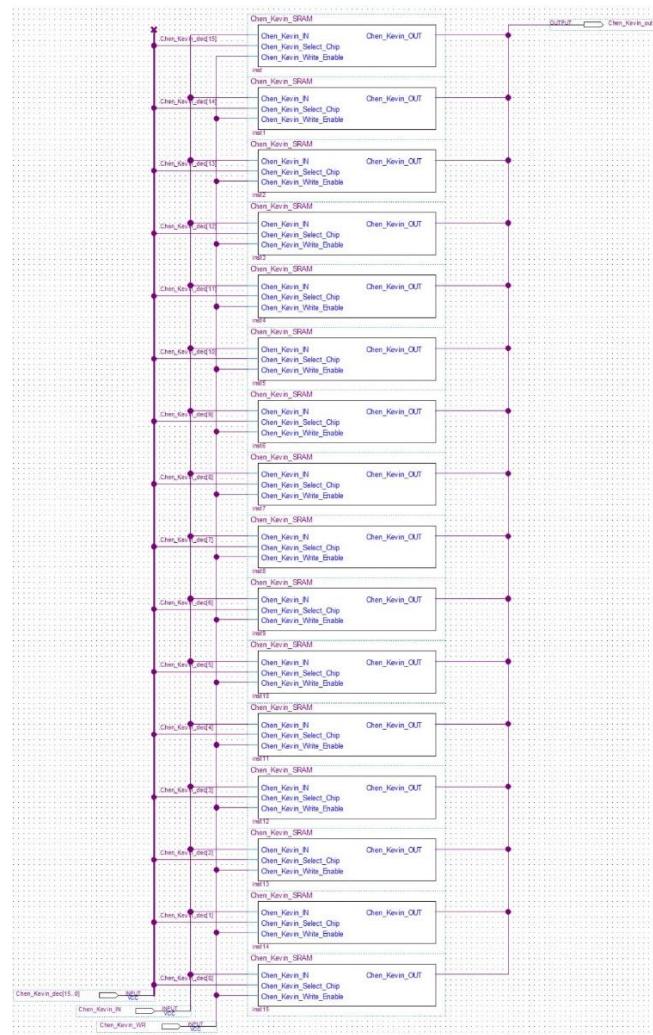
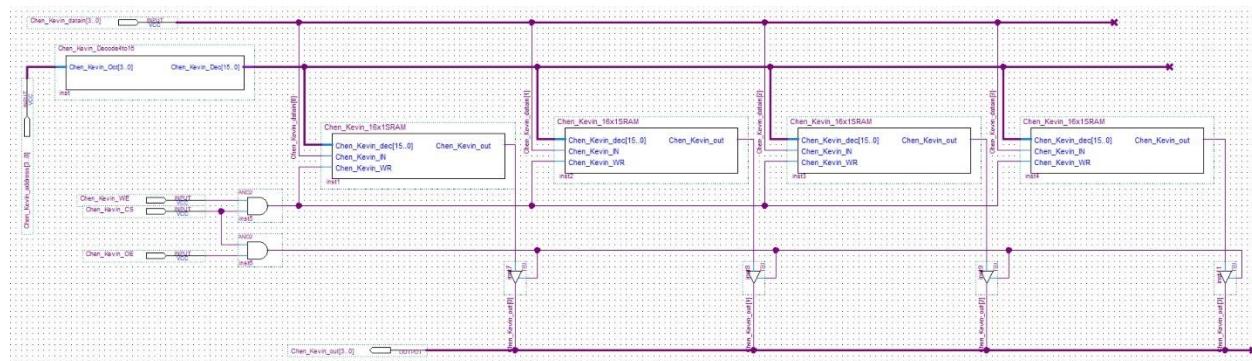


Figure 7: Block Diagram/Schematic File of 16 x 1 SRAM (Chen\_Kevin\_16x1SRAM.BDF)

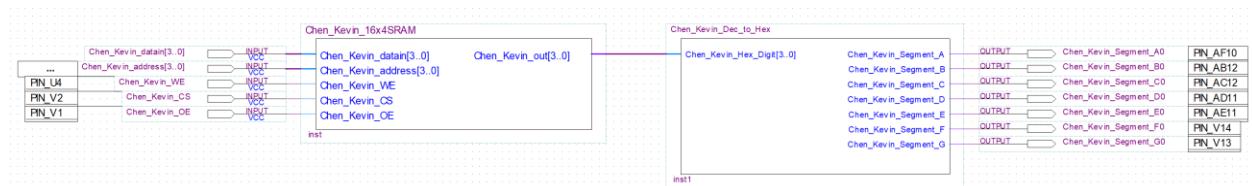
### 16 x 4 SRAM:

A  $16 \times 4$  SRAM is a circuit that's made of four  $16 \times 1$  SRAM that's all stacked into one file. It's also made from a 4-to-16 Decoder and a Decimal to Hexadecimal Translation. A 4-to-16 Decoder is a circuit that compares four inputs and lets out a result of sixteen outputs. The outputs are determined based on the inputs, and usually only one of the sixteen outputs are set to 1 and the others are set to 0, otherwise it's all set to 0. A Decimal to Hexadecimal Translation is a circuit that compares four inputs and lets out a result of seven outputs. The outputs are based on which seven-segments makes up the number corresponding to the inputs. The seven outputs would be used to display on a 7-segment-display, each output representing each segment.

The circuit is shown below:



**Figure 8: Block Diagram/Schematic File of 16 x 4 SRAM (Chen\_Kevin\_16x4SRAM.BDF)**



*Figure 9: Block Diagram/Schematic File of Final 16 x 4 SRAM*

(Chen Kevin 16x4SRAMFinal.BDF)

### 16 x 32 SRAM:

A 16 x 32 SRAM is a circuit that's made of eight 16 x 4 SRAM that's all stacked into one file.

It's also made from 8 Decimal to Hexadecimal Translation and a 4 x 8 Switches and Keys. A 4 x 8 Switches and Keys is a circuit that compares 12 inputs and lets out a result of 32 outputs.

Because we don't expect to have a pin for each 32 inputs, for the datain, this was designed to simplify the 32 inputs by using four keys and eight switches to represent each eight of the 32 inputs. For example, key[0] represents datain from 0 to 7, key[1] represents datain from 8 to 15, key[3] represents datain from 16 to 23, and key[4] represents datain from 24 to 31. The switches are from 0-7, which are used to configure each of the keys.

The circuit is shown below:

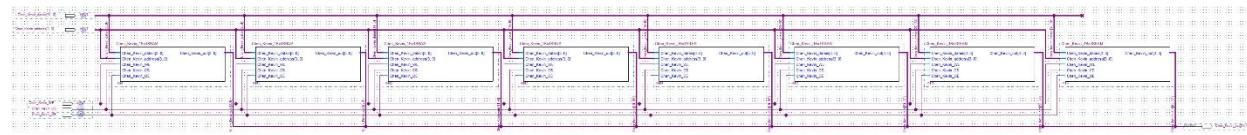


Figure 10: Block Diagram/Schematic File of 16 x 32 SRAM (Chen\_Kevin\_16x32SRAM.BDF)

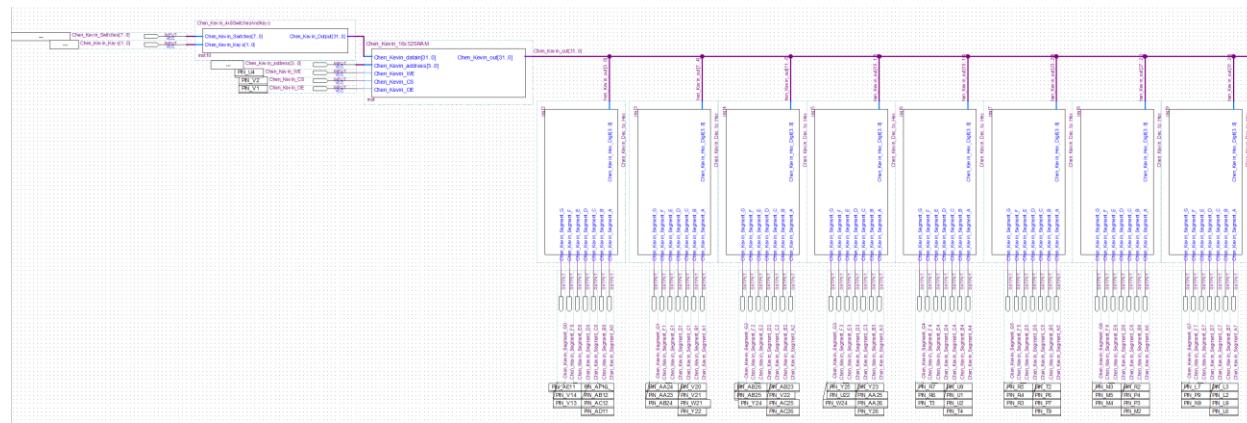


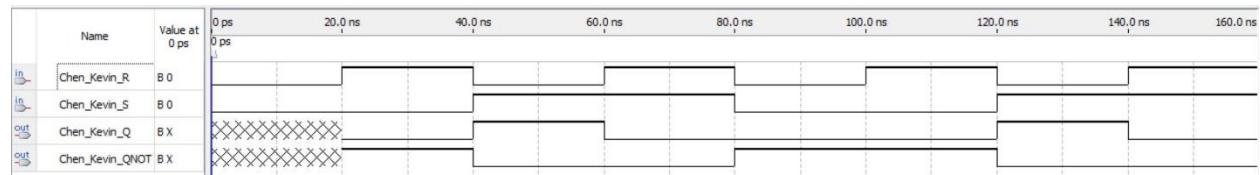
Figure 11: Block Diagram/Schematic File of Final 16 x 32 SRAM

(Chen\_Kevin\_16x32SRAMFinal.BDF)

**Simulations:****SR-Latch:**

The functionality/behavior of the SR-Latch is to follow the truth table (Table 1), which describes exactly what the output is supposed to be given a certain input.

A waveform was created using the block diagram/schematic file for SR-Latch (Figure 1) and it was given inputs to test all possible combinations. The outputs that it generated matches those from table 1, therefore, confirming the correctness of the block diagram/schematic file.



*Figure 12: Waveform Simulation of Block Diagram/Schematic File of SR-Latch*

*(Chen\_Kevin\_SR-LatchWaveform.VWF)*

Control SR-Latch:

The functionality/behavior of the Control SR-Latch is to follow the truth table (Table 2), which describes exactly what the output is supposed to be given a certain input.

A waveform was created using the block diagram/schematic file for Control SR-Latch (Figure 2) and it was given inputs to test all possible combinations. The outputs that it generated matches those from table 2, therefore, confirming the correctness of the block diagram/schematic file.

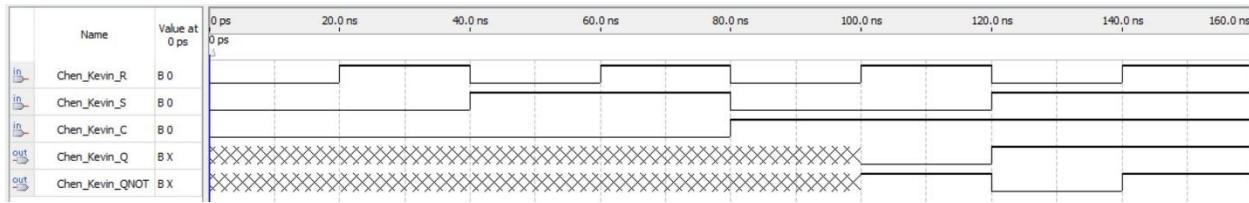


Figure 13: Waveform Simulation of Block Diagram/Schematic File of Control SR-Latch

(Chen\_Kevin\_Control\_SR-LatchWaveform.VWF)

### D-Latch:

The functionality/behavior of the D-Latch is to follow the truth table (Table 3), which describes exactly what the output is supposed to be given a certain input.

A waveform was created using the block diagram/schematic file for D-Latch (Figure 3) and it was given inputs to test all possible combinations. The outputs that it generated matches those from table 3, therefore, confirming the correctness of the block diagram/schematic file.

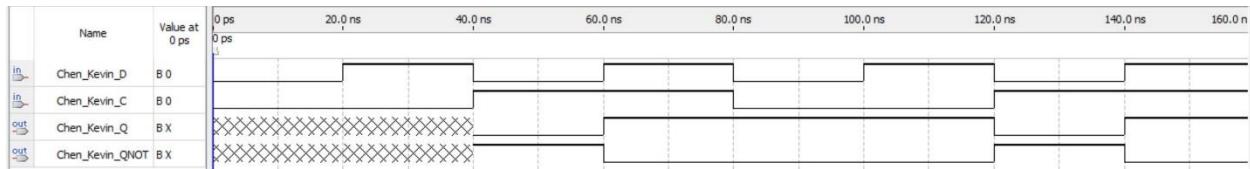


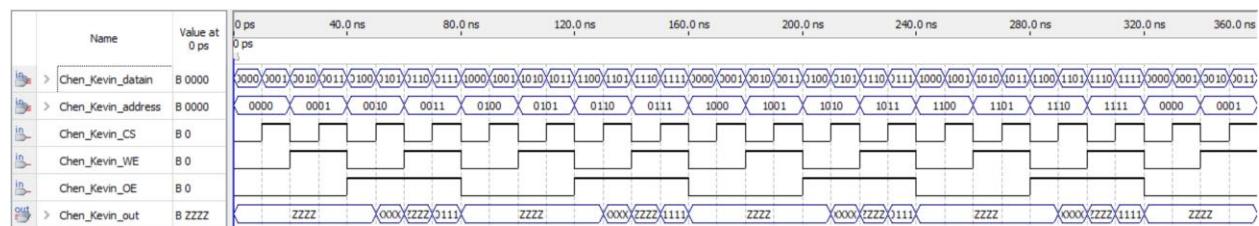
Figure 14: Waveform Simulation of Block Diagram/Schematic File of D-Latch

(Chen\_Kevin\_D-LatchWaveform.VWF)

## SRAM:

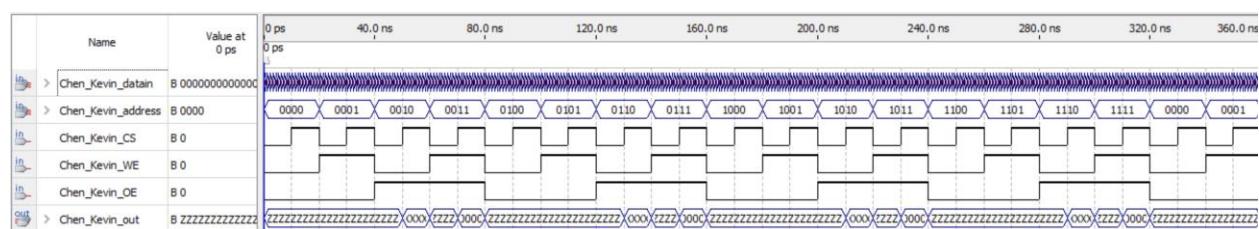
There are three functionality/behavior of the SRAM. One functionality/behavior serves to turn the chip on or off. The second functionality/behavior serves to write the “datain” input into the “address” input. The third functionality/behavior serves to display the output that was inputted.

A waveform was created using the block diagram/schematic file for SRAM (Figure 15 and 16) and it was given inputs to test all possible combinations. The output that it generates shows that the chip is turned on when “CS” is turned on; when “WE” is turned on, it allows to write; when “OE” is on, it displays the output.



*Figure 15: Waveform Simulation of Block Diagram/Schematic File of 16 x 4 SRAM*

(Chen\_Kevin\_16x4SRAMWaveform.VWF)



*Figure 16: Waveform Simulation of Block Diagram/Schematic File of 16 x32 SRAM*

(Chen\_Kevin\_16x32SRAMWaveform.VWF)

**Demonstration:****16 x 4 SRAM:**

In address 0000, I inputted the binary code 0011 and had it display on the DE2 board. After that, I swapped into address 0001, and inputted the binary code 1001 and had it display on the DE2 board.

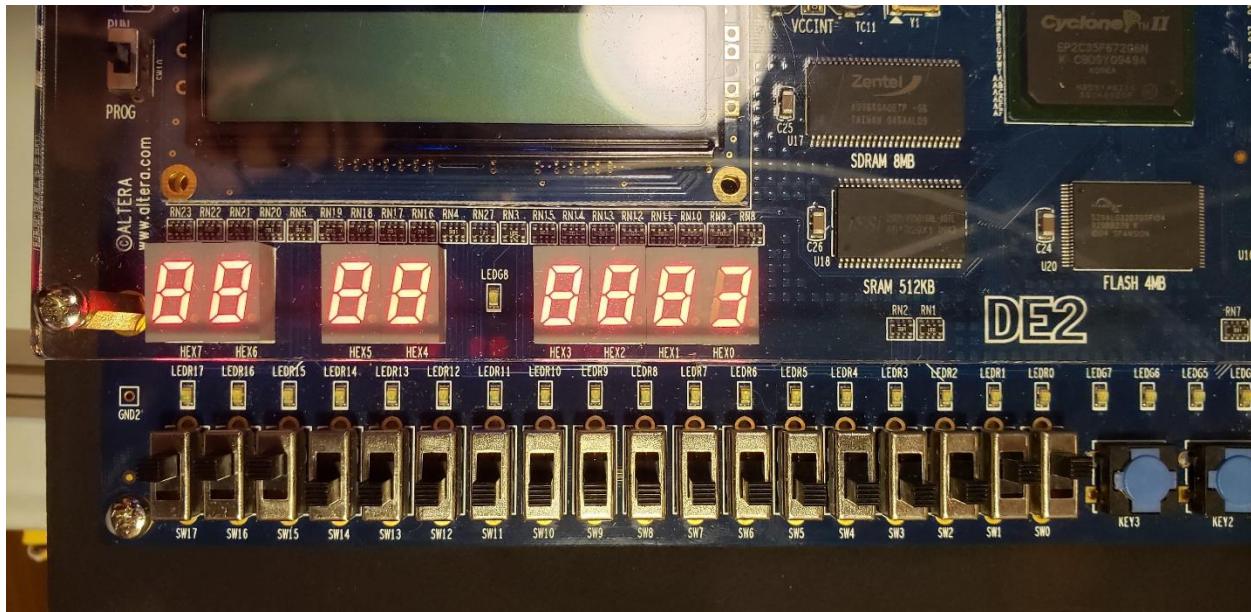


Figure 17: Input 3 Into Address 0 on DE2 Board

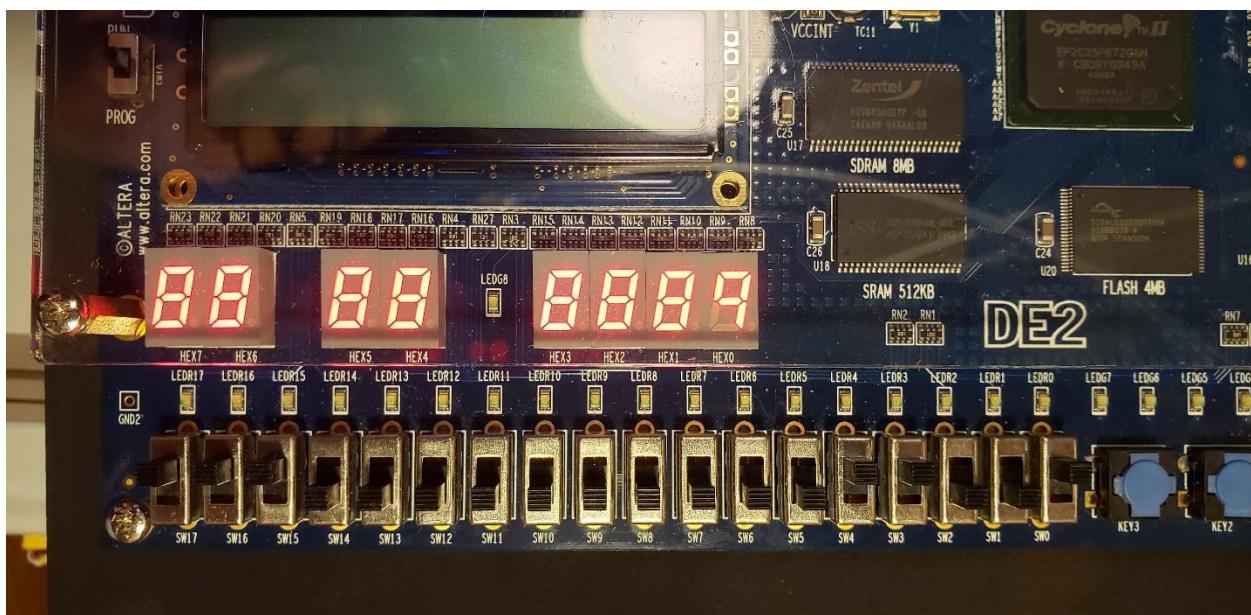


Figure 18: Input 9 Into Address 1 on DE2 Board

16 x 32 SRAM:

In address 0000, I inputted the binary code 1000 0111 0110 0101 0100 0011 0010 0001 and had it display on the DE2 board. After that, I swapped into address 0001, and inputted the binary code 1010 1011 1100 1101 1110 1111 0000 0001 and had it display on the DE2 board.

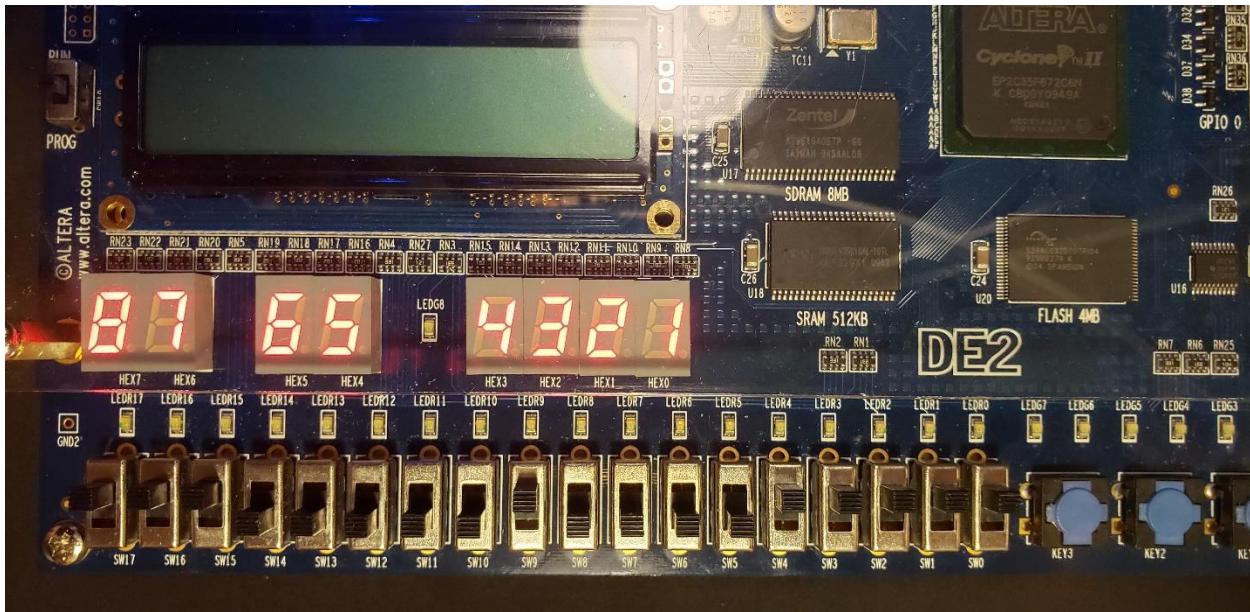


Figure 19: Input 8 7 6 5 4 3 2 1 Into Address 0 on DE2 Board

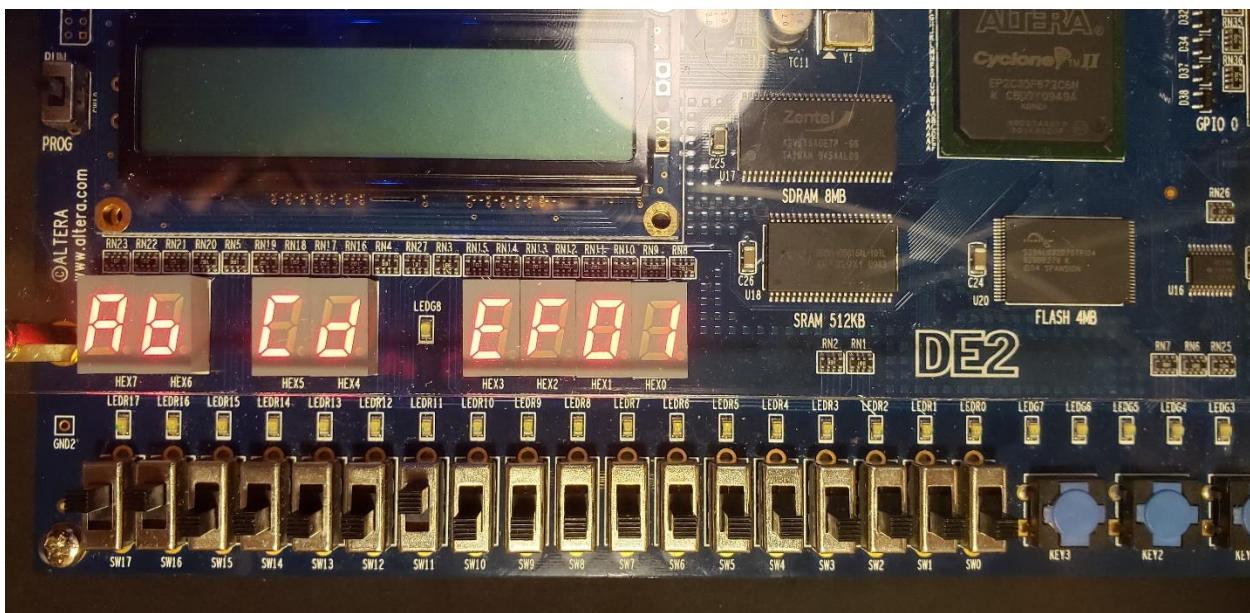


Figure 20: Input A b C d E F 0 1 Into Address 1 on DE2 Board

**Conclusion:**

This lab taught me how to create a static random-access memory chip using D-Latches. It also taught me how to allow the output of the SRAM to be connected to a 4-to-7 decoder that outputs to the 7-segment-display. I was able to view how a SR-Latch could become into a Control SR-Latch, then into a D-Latch, then into a Master-Slave D Flip-Flop, then into a SRAM.

1. What is the difference between a latch and a flip-flop?

The difference between a latch and a flip-flop is that a latch is level triggered so the output can change anytime based on the inputs, and a flip-flop is edge triggered so the output can only change when the clock changes.

2. Explain the behavior of a SR-Latch

The behavior of a SR-Latch is either give the same outputs as the previous output or to set an output to 1 or reset an output to 0.

3. Explain how the gated (control) SR-Latch is a modified SR-Latch

The gated (control) SR-Latch is a modified SR-Latch because it performs the same action as a SR-Latch except with an additional control input. When the control input is not active, then nothing changes, but when it is active, it acts the same as a SR-Latch.

4. Explain how the D-Latch is a modified gated SR-Latch

The D-Latch is a modified gated SR-Latch by how it uses the same control input and combines the S and R input into one input, called “D”. The “D” input acts the same way as a SR-Latch except now you can’t have both set and reset; it’s either or.

5. Explain how to build a Master-Slave flip-flop.

A Master-Slave Flip-Flop is built by connecting two latches in series. Figure 4 shows how it's built.

6. Explain how to build SRAM cells from flip-flops.

A SRAM cell is built by using a Master-Slave D Flip-Flop and connecting it to a few logic gates and a tristate logic buffer. Figure 6 shows how it's built.

7. Why do we want to avoid the scenario when S=1 and R=1 in an SR-Latch? Could this scenario happen in a D-Latch? Why or why not?

We want to avoid the scenario when S=1 and R=1 in an SR-Latch because we can't have a device both set and reset at the same time. This scenario can't happen in a D-Latch because a D-Latch combines the set and reset inputs into one input, which makes it so it can either be set or reset.

8. What is the difference between edge-triggered and level triggered devices? Is the flip-flop used for the SRAM cell in this lab level or edge triggered?

The difference between edge-triggered and level triggered devices is that edge-triggered devices depend on the clock signal for a change in output, while level triggered devices depend on a change in inputs for a change in output. The flip-flop used for the SRAM cell in this lab is edge triggered.

9. Can you explain how an SRAM works?

An SRAM works by first enabling the “Chip\_Select”, which turns the SRAM on. Then by giving it a specific input and a specific address, where the input will be stored in, the input would be written into the address when the “Write\_Enable” is enabled. By enabling the “Output\_Enable”, it's able to display what was written into the address.

## Appendix:

### 4-to-16 Decoder:

```
1 library IEEE;
2 use IEEE.STD_LOGIC_1164.all;
3
4 entity Chen_Kevin_Decode4to16 is
5 port(
6     Chen_Kevin_Oct : in std_logic_vector(3 downto 0);
7     Chen_Kevin_Dec : out std_logic_vector(15 downto 0));
8 end Chen_Kevin_Decode4to16;
9
10 architecture Arch of Chen_Kevin_Decode4to16 is
11
12 begin
13
14 with Chen_Kevin_Oct select
15     Chen_Kevin_Dec <=    "0000000000000001" when "0000",
16                             "0000000000000010" when "0001",
17                             "00000000000000100" when "0010",
18                             "000000000000001000" when "0011",
19                             "0000000000000010000" when "0100",
20                             "00000000000000100000" when "0101",
21                             "00000000001000000" when "0110",
22                             "000000000100000000" when "0111",
23                             "0000000010000000000" when "1000",
24                             "0000000100000000000" when "1001",
25                             "0000010000000000000" when "1010",
26                             "0000100000000000000" when "1011",
27                             "0001000000000000000" when "1100",
28                             "0010000000000000000" when "1101",
29                             "0100000000000000000" when "1110",
30                             "1000000000000000000" when "1111",
31                             "0000000000000000000" when others;
32 end arch;
```

Figure 21: VHDL Code for 4-to-16 Decoder (Chen\_Kevin\_Decode4to16.VHD)

## Decimal to Hexadecimal Translation:

```

1 LIBRARY IEEE;
2 USE IEEE.STD.LOGIC_1164.all;
3 USE IEEE.STD.LOGIC_UNSIGNED.all;
4 USE IEEE.STD.LOGIC_UNSIGNED.all;
5 --Hexadecimal to 7 Segment Decoder for LED Display
6
7 ENTITY Chen_Kevin_Dec_to_Hex IS
8 PORT ( Chen_Kevin_Hex_Digit : IN STD LOGIC VECTOR (3 DOWNTO 0);
9 Chen_Kevin_Segment_A, Chen_Kevin_Segment_B, Chen_Kevin_Segment_C, Chen_Kevin_Segment_D, Chen_Kevin_Segment_E, Chen_Kevin_Segment_F, Chen_Kevin_Segment_G : OUT std_logic);
10 END Chen_Kevin_Dec_to_Hex;
11
12 ARCHITECTURE Arc OF Chen_Kevin_Dec_to_Hex IS
13 SIGNAL Chen_Kevin_Segment_Data : STD_LOGIC_VECTOR (6 DOWNTO 0);
14 BEGIN
15 PROCESS (Chen_Kevin_Hex_Digit)
16 --HEX to 7 Segment Decoder for LED Display
17 BEGIN -- Hex-digit is the four bit binary value to display
18
19 CASE Chen_Kevin_Hex_Digit IS
20 WHEN "0000" =>
21 Chen_Kevin_Segment_Data <= "1111110";
22 WHEN "0001" =>
23 Chen_Kevin_Segment_Data <= "0110000";
24 WHEN "0010" =>
25 Chen_Kevin_Segment_Data <= "1101101";
26 WHEN "0011" =>
27 Chen_Kevin_Segment_Data <= "1111001";
28 WHEN "0100" =>
29 Chen_Kevin_Segment_Data <= "0110011";
29 WHEN "0101" =>
30 Chen_Kevin_Segment_Data <= "1011101";
31 WHEN "0110" =>
32 Chen_Kevin_Segment_Data <= "1011111";
33 WHEN "0111" =>
34 Chen_Kevin_Segment_Data <= "1110000";
35 WHEN "1000" =>
36 Chen_Kevin_Segment_Data <= "1111111";
37 WHEN "1001" =>
38 Chen_Kevin_Segment_Data <= "1110011";
39 WHEN "1010" =>
40 Chen_Kevin_Segment_Data <= "1110111";
41 WHEN "1011" =>
42 Chen_Kevin_Segment_Data <= "0011111";
43 WHEN "1100" =>
44 Chen_Kevin_Segment_Data <= "1001110";
45 WHEN "1101" =>
46 Chen_Kevin_Segment_Data <= "0111101";
47 WHEN "1110" =>
48 Chen_Kevin_Segment_Data <= "1001111";
49 WHEN "1111" =>
50 Chen_Kevin_Segment_Data <= "1000111";
51 END CASE;
52 END PROCESS;
53 -- extract segment data bits and invert
54 -- LED driver circuit is inverted
55 Chen_Kevin_Segment_A <= NOT Chen_Kevin_Segment_Data (0);
56 Chen_Kevin_Segment_B <= NOT Chen_Kevin_Segment_Data (1);
57 Chen_Kevin_Segment_C <= NOT Chen_Kevin_Segment_Data (2);
58 Chen_Kevin_Segment_D <= NOT Chen_Kevin_Segment_Data (3);
59 Chen_Kevin_Segment_E <= NOT Chen_Kevin_Segment_Data (4);
60 Chen_Kevin_Segment_F <= NOT Chen_Kevin_Segment_Data (5);
61 Chen_Kevin_Segment_G <= NOT Chen_Kevin_Segment_Data (6);
62 END Arc;

```

Figure 22: VHDL Code for Decimal to Hexadecimal Translation

(Chen\_Kevin\_Dec\_to\_Hex.VHD)

16 x 4 SRAM Pin Assignments:

<b>Variable Name</b>	<b>Variable Type</b>	<b>Signal Name</b>	<b>FPGA Pin No.</b>
Chen_Kevin_datain[0]	Input	SW[0]	PIN_N25
Chen_Kevin_datain[1]	Input	SW[1]	PIN_N26
Chen_Kevin_datain[2]	Input	SW[2]	PIN_P25
Chen_Kevin_datain[3]	Input	SW[3]	PIN_AE14
Chen_Kevin_address[0]	Input	SW[4]	PIN_AF14
Chen_Kevin_address[1]	Input	SW[5]	PIN_AD13
Chen_Kevin_address[2]	Input	SW[6]	PIN_AC13
Chen_Kevin_address[3]	Input	SW[7]	PIN_C13
Chen_Kevin_WE	Input	SW[15]	PIN_U4
Chen_Kevin_OE	Input	SW[16]	PIN_V1
Chen_Kevin_CS	Input	SW[17]	PIN_V2
Chen_Kevin_Segment_A0	Output	HEX0[0]	PIN_AF10
Chen_Kevin_Segment_B0	Output	HEX0[1]	PIN_AB12
Chen_Kevin_Segment_C0	Output	HEX0[2]	PIN_AC12
Chen_Kevin_Segment_D0	Output	HEX0[3]	PIN_AD11
Chen_Kevin_Segment_E0	Output	HEX0[4]	PIN_AE11
Chen_Kevin_Segment_F0	Output	HEX0[5]	PIN_V14
Chen_Kevin_Segment_G0	Output	HEX0[6]	PIN_V13

Table 6: Pin Assignments for 16 x 4 SRAM

## 16 x 4 SRAM Pin Planner:

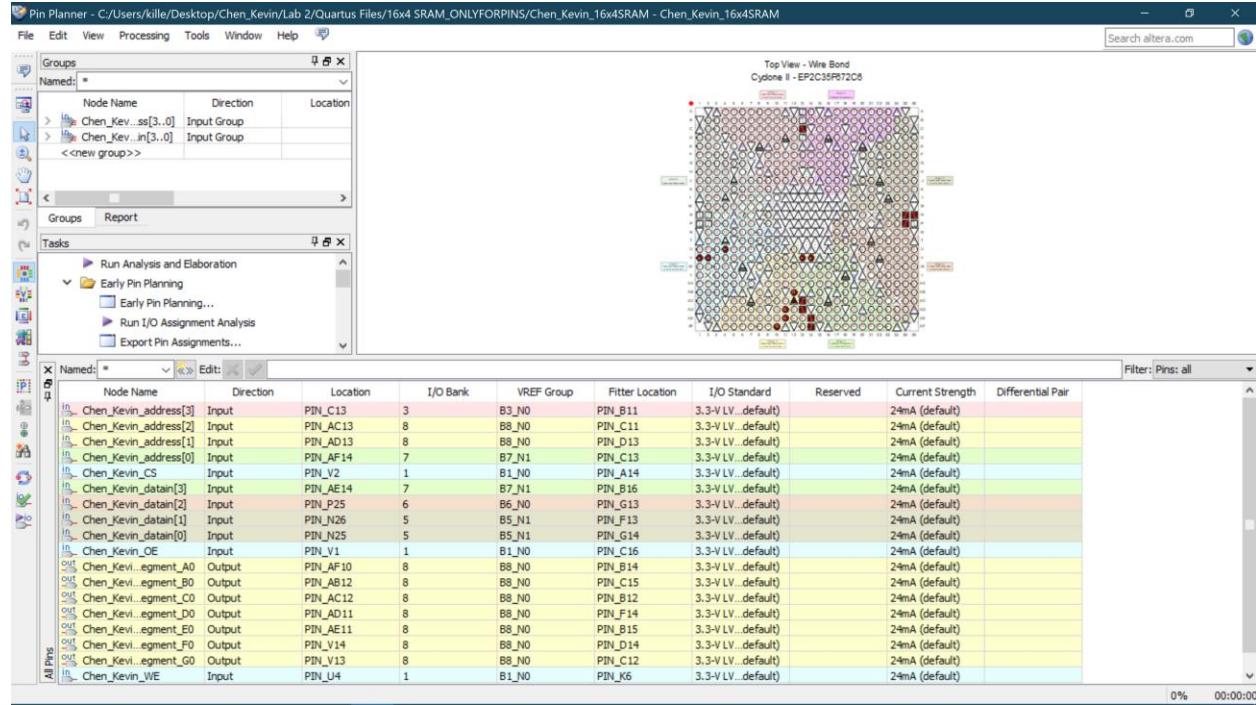


Figure 23: Pin Planner of Block Diagram/Schematic File of 16 x 4 SRAM

4 x 8 Switches and Keys:

```
1 library IEEE;
2 use IEEE.STD_LOGIC_1164.ALL;
3
4 entity Chen_Kevin_4x8SwitchesAndKeys is
5   port(
6     Chen_Kevin_Switches : in std_logic_vector(7 downto 0);
7     Chen_Kevin_Keys : in std_logic_vector(1 downto 0);
8     Chen_Kevin_Output : out std_logic_vector(31 downto 0));
9 end Chen_Kevin_4x8SwitchesAndKeys;
10
11 architecture Arch of Chen_Kevin_4x8SwitchesAndKeys is
12
13 begin
14   process(Chen_Kevin_Keys)
15   begin
16     case Chen_Kevin_Keys is
17       WHEN "00" =>
18         Chen_Kevin_Output(7 downto 0) <= Chen_Kevin_Switches;
19       WHEN "01" =>
20         Chen_Kevin_Output(15 downto 8) <= Chen_Kevin_Switches;
21       WHEN "10" =>
22         Chen_Kevin_Output(23 downto 16) <= Chen_Kevin_Switches;
23       WHEN "11" =>
24         Chen_Kevin_Output(31 downto 24) <= Chen_Kevin_Switches;
25     end case;
26   end process;
27 end Arch;
```

Figure 24: VHDL Code for 4 x 8 Switches and Keys (Chen\_Kevin\_4x8SwitchesAndKeys.VHD)

16 x 32 SRAM Pin Assignments:

<b>Variable Name</b>	<b>Variable Type</b>	<b>Signal Name</b>	<b>FPGA Pin No.</b>
Chen_Kevin_Keys[0]	Input	SW[0]	PIN_N25
Chen_Kevin_Keys[1]	Input	SW[1]	PIN_N26
Chen_Kevin_Switches[0]	Input	SW[2]	PIN_P25
Chen_Kevin_Switches[1]	Input	SW[3]	PIN_AE14
Chen_Kevin_Switches[2]	Input	SW[4]	PIN_AF14
Chen_Kevin_Switches[3]	Input	SW[5]	PIN_AD13
Chen_Kevin_Switches[4]	Input	SW[6]	PIN_AC13
Chen_Kevin_Switches[5]	Input	SW[7]	PIN_C13
Chen_Kevin_Switches[6]	Input	SW[8]	PIN_B13
Chen_Kevin_Switches[7]	Input	SW[9]	PIN_A13
Chen_Kevin_address[0]	Input	SW[10]	PIN_N1
Chen_Kevin_address[1]	Input	SW[11]	PIN_P1
Chen_Kevin_address[2]	Input	SW[12]	PIN_P2
Chen_Kevin_address[3]	Input	SW[13]	PIN_T7
Chen_Kevin_WE	Input	SW[15]	PIN_U4
Chen_Kevin_OE	Input	SW[16]	PIN_V1
Chen_Kevin_CS	Input	SW[17]	PIN_V2

Table 7a: Pin Assignments for 16 x 32 SRAM

Variable Name	Variable Type	Signal Name	FPGA Pin No.
Chen_Kevin_Segment_A0	Output	HEX0[0]	PIN_AF10
Chen_Kevin_Segment_B0	Output	HEX0[1]	PIN_AB12
Chen_Kevin_Segment_C0	Output	HEX0[2]	PIN_AC12
Chen_Kevin_Segment_D0	Output	HEX0[3]	PIN_AD11
Chen_Kevin_Segment_E0	Output	HEX0[4]	PIN_AE11
Chen_Kevin_Segment_F0	Output	HEX0[5]	PIN_V14
Chen_Kevin_Segment_G0	Output	HEX0[6]	PIN_V13
Chen_Kevin_Segment_A1	Output	HEX1[0]	PIN_V20
Chen_Kevin_Segment_B1	Output	HEX1[1]	PIN_V21
Chen_Kevin_Segment_C1	Output	HEX1[2]	PIN_W21
Chen_Kevin_Segment_D1	Output	HEX1[3]	PIN_Y22
Chen_Kevin_Segment_E1	Output	HEX1[4]	PIN_AA24
Chen_Kevin_Segment_F1	Output	HEX1[5]	PIN_AA23
Chen_Kevin_Segment_G1	Output	HEX1[6]	PIN_AB24
Chen_Kevin_Segment_A2	Output	HEX2[0]	PIN_AB23
Chen_Kevin_Segment_B2	Output	HEX2[1]	PIN_V22
Chen_Kevin_Segment_C2	Output	HEX2[2]	PIN_AC25
Chen_Kevin_Segment_D2	Output	HEX2[3]	PIN_AC26
Chen_Kevin_Segment_E2	Output	HEX2[4]	PIN_AB26
Chen_Kevin_Segment_F2	Output	HEX2[5]	PIN_AB25
Chen_Kevin_Segment_G2	Output	HEX2[6]	PIN_Y24

Table 7b: Pin Assignments for 16 x 32 SRAM

Variable Name	Variable Type	Signal Name	FPGA Pin No.
Chen_Kevin_Segment_A3	Output	HEX3[0]	PIN_Y23
Chen_Kevin_Segment_B3	Output	HEX3[1]	PIN_AA25
Chen_Kevin_Segment_C3	Output	HEX3[2]	PIN_AA26
Chen_Kevin_Segment_D3	Output	HEX3[3]	PIN_Y26
Chen_Kevin_Segment_E3	Output	HEX3[4]	PIN_Y25
Chen_Kevin_Segment_F3	Output	HEX3[5]	PIN_U22
Chen_Kevin_Segment_G3	Output	HEX3[6]	PIN_W24
Chen_Kevin_Segment_A4	Output	HEX4[0]	PIN_U9
Chen_Kevin_Segment_B4	Output	HEX4[1]	PIN_U1
Chen_Kevin_Segment_C4	Output	HEX4[2]	PIN_U2
Chen_Kevin_Segment_D4	Output	HEX4[3]	PIN_T4
Chen_Kevin_Segment_E4	Output	HEX4[4]	PIN_R7
Chen_Kevin_Segment_F4	Output	HEX4[5]	PIN_R6
Chen_Kevin_Segment_G4	Output	HEX4[6]	PIN_T3
Chen_Kevin_Segment_A5	Output	HEX5[0]	PIN_T2
Chen_Kevin_Segment_B5	Output	HEX5[1]	PIN_P6
Chen_Kevin_Segment_C5	Output	HEX5[2]	PIN_P7
Chen_Kevin_Segment_D5	Output	HEX5[3]	PIN_T9
Chen_Kevin_Segment_E5	Output	HEX5[4]	PIN_R5
Chen_Kevin_Segment_F5	Output	HEX5[5]	PIN_R4
Chen_Kevin_Segment_G5	Output	HEX5[6]	PIN_R3

Table 7c: Pin Assignments for 16 x 32 SRAM

Variable Name	Variable Type	Signal Name	FPGA Pin No.
Chen_Kevin_Segment_A6	Output	HEX6[0]	PIN_R2
Chen_Kevin_Segment_B6	Output	HEX6[1]	PIN_P4
Chen_Kevin_Segment_C6	Output	HEX6[2]	PIN_P3
Chen_Kevin_Segment_D6	Output	HEX6[3]	PIN_M2
Chen_Kevin_Segment_E6	Output	HEX6[4]	PIN_M3
Chen_Kevin_Segment_F6	Output	HEX6[5]	PIN_M5
Chen_Kevin_Segment_G6	Output	HEX6[6]	PIN_M4
Chen_Kevin_Segment_A7	Output	HEX7[0]	PIN_L3
Chen_Kevin_Segment_B7	Output	HEX7[1]	PIN_L2
Chen_Kevin_Segment_C7	Output	HEX7[2]	PIN_L9
Chen_Kevin_Segment_D7	Output	HEX7[3]	PIN_L6
Chen_Kevin_Segment_E7	Output	HEX7[4]	PIN_L7
Chen_Kevin_Segment_F7	Output	HEX7[5]	PIN_P9
Chen_Kevin_Segment_G7	Output	HEX7[6]	PIN_N9

Table 7d: Pin Assignments for 16 x 32 SRAM

## 16 x 32 SRAM Pin Planner:

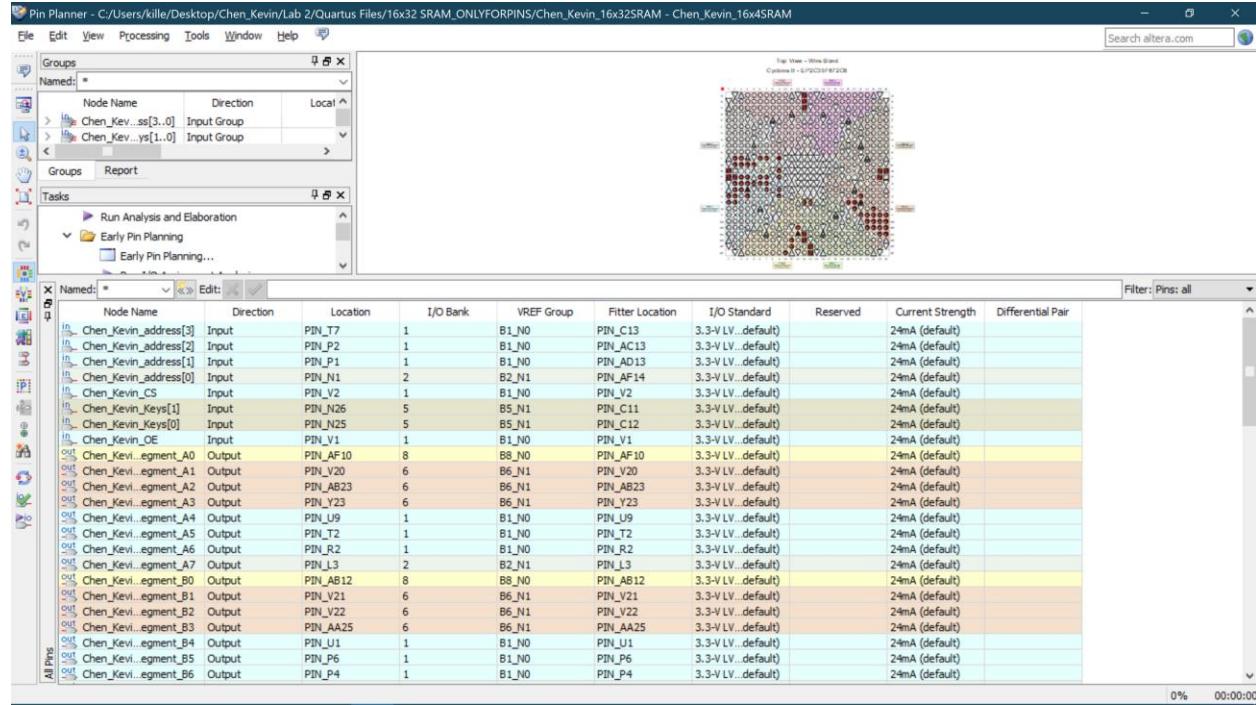


Figure 25a: Pin Planner of Block Diagram/Schematic File of 16 x 32 SRAM

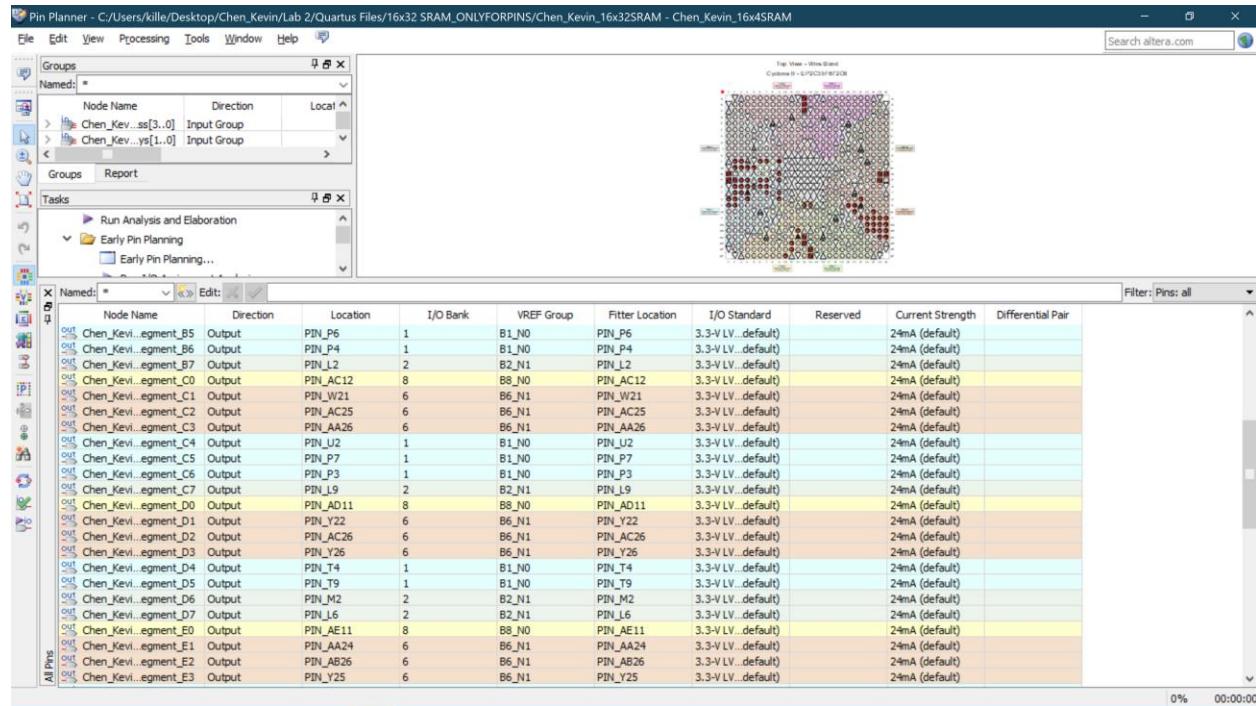


Figure 25b: Pin Planner of Block Diagram/Schematic File of 16 x 32 SRAM

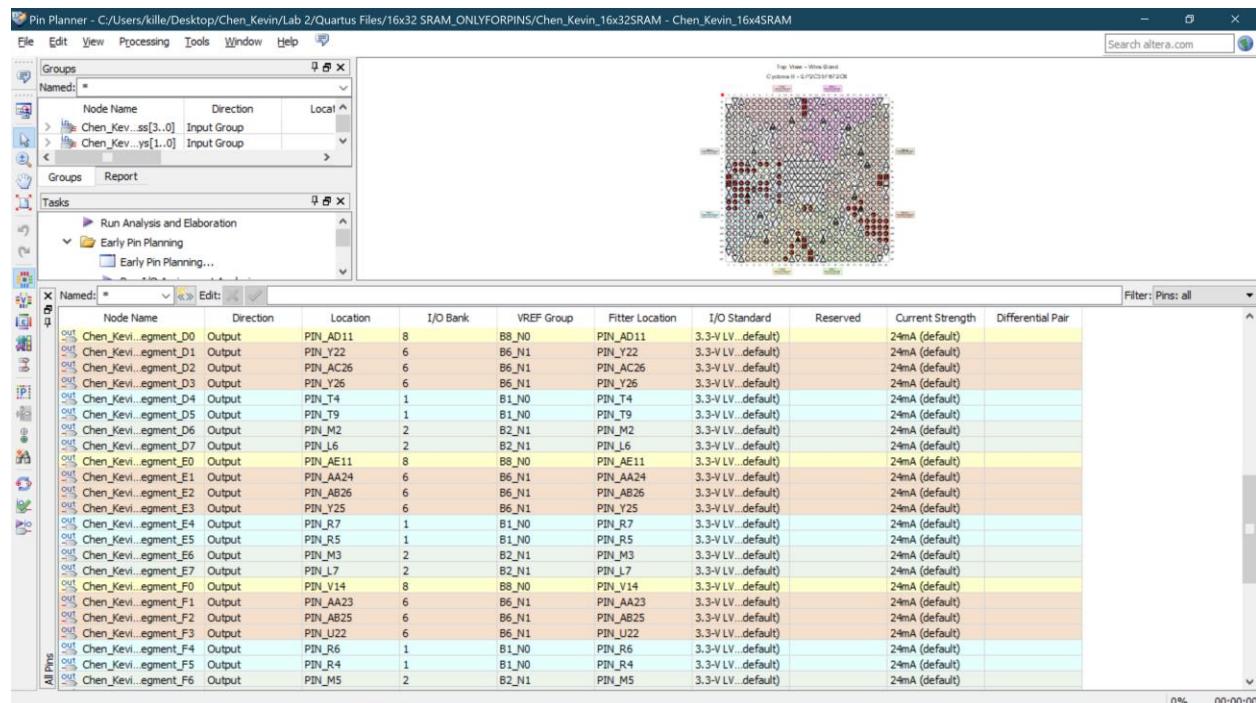


Figure 25c: Pin Planner of Block Diagram/Schematic File of 16 x 32 SRAM

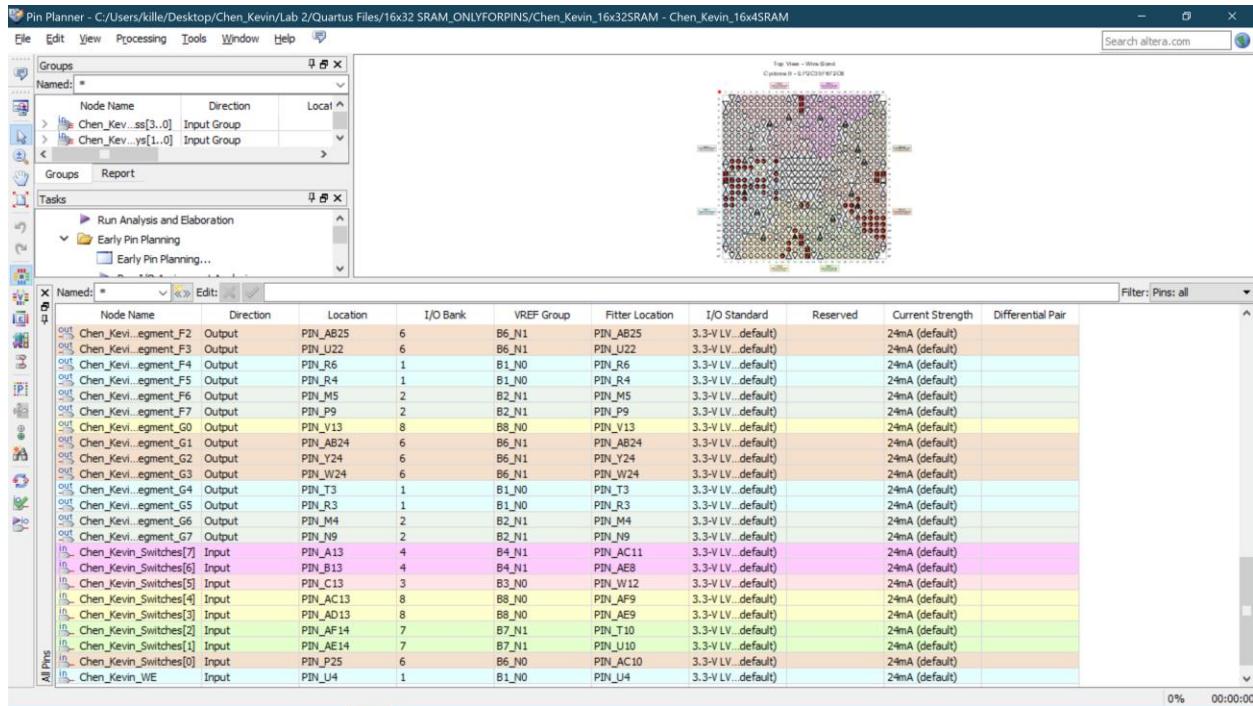


Figure 25d: Pin Planner of Block Diagram/Schematic File of 16 x 32 SRAM